

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE



In re PATENT APPLICATION of
Inventor(s): SOININEN et al.

Appln. No.: 09
Series Code ↑ Serial No. ↑

Group Art Unit: Not Yet Assigned

Filed: August 28, 2001

Examiner: Not Yet Assigned

Title: IP ROUTING OPTIMIZATION IN AN ACCESS NETWORK Atty. Dkt. P 282821 2990421US/A/HER
M# Client Ref

Date: August 29, 2001

SUBMISSION OF PRIORITY DOCUMENT IN ACCORDANCE WITH THE REQUIREMENTS OF RULE 55

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<u>Application No.</u>	<u>Country of Origin</u>	<u>Filed</u>
990510	FINLAND	March 9, 1999
991260	FINLAND	June 2, 1999

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Helsinki 10.8.2001

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PRIORITY DOCUMENT

11000 U.S. PRO
09/940577
08/29/01

Hakija
Applicant

Nokia Telecommunications Oy
Helsinki

Patenttihakemus nro
Patent application no

990510

Tekemispäivä
Filing date

09.03.1999

Kansainvälinen luokka
International class

H04Q

Keksinnön nimitys
Title of invention

"IP routing optimization in a packet radio network"
(IP-reitityksen optimointi pakettiradioverkossa)

Hakijan nimi on hakemusdiaariin 05.12.1999 tehdyn nimenmuutoksen jälkeen Nokia Networks Oy.

The application has according to an entry made in the register of patent applications on 05.12.1999 with the name changed into Nokia Networks Oy.

Täten todistetaan, että oheiset asiakirjat ovat tarkkoja jäljennöksiä patentti- ja rekisterihallitukselle alkuaan annetuista selityksestä, patenttivaatimuksista, tiivistelmästä ja piirustuksista.

This is to certify that the annexed documents are true copies of the description, claims, abstract and drawings originally filed with the Finnish Patent Office.


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IP routing optimization in a packet radio network

Field of the Invention

The invention relates to a mechanism for optimizing the routing of an Internet-type protocol traffic in a packet radio network.

5 Background of the Invention

Mobile communications system refers generally to any telecommunications system which enable a wireless communication when users are moving within the service area of the system. A typical mobile communications system is a Public Land Mobile Network (PLMN). Often the mobile communications network is an access network providing a user with a wireless access to
10 external networks, hosts, or services offered by specific service providers.

The general packet radio service GPRS is a new service in the GSM system (Global System for Mobile communication). A subnetwork comprises a number of packet data service nodes SN, which in this application will
15 be referred to as serving GPRS support nodes SGSN. Each SGSN is connected to the GSM mobile communication network (typically to a base station controller BSC or a base station BTS in a base station system) so that the SGSN can provide a packet service for mobile data terminals via several base stations, i.e. cells. The intermediate mobile communication network provides a
20 radio access and a packet-switched data transmission between the SGSN and mobile data terminals. Different subnetworks are in turn connected to an external data network, e.g. to a public switched data network PSPDN, via GPRS gateway support nodes GGSN. The GPRS service thus allows to provide packet data transmission between mobile data terminals and external
25 data networks when the GSM network functions as a radio access network RAN.

Third generation mobile systems, such as Universal Mobile Communications system (UMTS) and Future Public Land Mobile Telecommunications system (FPLMTS), later renamed as IMT-2000 (International Mobile
30 Telecommunication 2000), are being developed. In the UMTS architecture a UMTS terrestrial radio access network, UTRAN, consists of a set of radio access networks RAN (also called radio network subsystem RNS) connected to the core network (CN). Each RAN is responsible for the resources of its set of cells. For each connection between a mobile station MS and the UTRAN, one
35 RAN is a serving RAN. A RAN consists of a radio network controller RNC and

a multiplicity of base stations BS. One core network which will use the UMTS radio access network is the GPRS.

One of the main targets in the development of mobile communication network is to provide IP (Internet Protocol) service with a standard IP backbone which would use a combination of mobile network mobility management in the mobile networks and Mobile IP. The basic IP concept does not support the mobility of the user: the IP addresses are assigned to network interfaces in dependence on their physical location. In fact, the first field of an IP address (the NETID) is common to all interfaces that are linked to the same Internet subnet. This scheme prevents the user (the mobile host) from keeping its address while moving over different Internet subnets, i.e. while changing the physical interface.

In order to enhance the mobility in the Internet a Mobile IP protocol for IP version 4 have been introduced by the Internet Engineering Task Force (IETF) in the standard RFC2002. Mobile IP enables the routing of IP datagrams to mobile hosts, independently of the point of attachment in the sub-network. The mobile IP protocol introduces following new functional or architectural entities.

'Mobile Node MN' (also called Mobile Host MH) refers to a host that changes its point of attachment from one network or subnetwork to another. A mobile node may change its location without changing its IP address; it may continue to communicate with other Internet nodes at any location using its (constant) IP address. 'Mobile Station (MS)' is a mobile node having a radio interface to the network. A 'Tunnel' is the path followed by a datagram when it is encapsulated. The model is that, while it is encapsulated, a datagram is routed to a known decapsulation agent, which decapsulates the datagram and then correctly delivers it to its ultimate destination. Each mobile node is connected to a home agent over a unique tunnel, identified by a tunnel identifier which is unique to a given Foreign Agent/Home Agent pair.

'Home Network' is the IP network to which a user logically belongs. Physically, it can be e.g. a local area network (LAN) connected via a router to the Internet. 'Home Address' is an address that is assigned to a mobile node for an extended period of time. It may remain unchanged regardless of where the MN is attached to the Internet. Alternatively, it could be assigned from a pool of addresses.

'Mobility Agent' is either a home agent or a foreign agent. 'Home Agent HA' is a routing entity on a mobile node's home network which tunnels packets for delivery to the mobile node when it is away from home, and maintains current location information for the mobile node. It tunnels datagrams for delivery to, and, optionally, detunnels datagrams from, a mobile node when the mobile node is away from home. 'Foreign Agent FA' refers to a routing entity in a mobile node's visited network which provides routing services to the mobile node while registered, thus allowing a mobile node to utilise its home network address. The foreign agent detunnels and delivers packets to the mobile node that were tunnelled by the mobile node's home agent. For datagrams sent by a mobile node, the foreign agent may serve as a default router for registered mobile nodes.

RFC2002 defines 'Care-of Address' (COA) as the termination point of a tunnel toward a mobile node, for datagrams forwarded to the mobile node while it is away from home. The protocol can use two different types of care-of address: a "foreign agent care-of address" is an address announced by a foreign agent with which the mobile node is registered, and a "co-located care-of address" is an externally obtained local address which the mobile node has acquired in the network. An MN may have several COAs at the same time. An MN's COA is registered with its HA. The list of COAs is updated when the mobile node receives advertisements from foreign agents. If an advertisement expires, its entry or entries should be deleted from the list. One foreign agent can provide more than one COA in its advertisements. 'Mobility Binding' is the association of a home address with a care-of address, along with the remaining lifetime of that association. An MN registers its COA with its HA by sending a Registration Request. The HA replies with a Registration Reply and retains a binding for the MN.

A single generic mobility handling mechanism that allows roaming between all types of access networks would allow the user to conveniently move between fixed and mobile networks, between public and private networks as well as between PLMN's with different access technologies. Therefore, mechanisms supporting the Mobile IP functionality are being developed also in mobile communication systems, such as UMTS and GPRS.

It is desired that the Mobile IP will be implemented as an overlay of the UMTS/GPRS network while maintaining backwards compatibility with present systems, assuming minimal modifications in the GPRS standards and on

networks whose operators do not want to support MIP. Fig. 1 illustrates the minimum configuration for a GPRS operator, who wishes to offer the mobile IP service. The current GPRS structure is kept and handles the mobility within the PLMN, while MIP allows user to roam between other systems, such as LAN's, and UMTS without losing an ongoing session. In Fig. 1 the foreign agents FA are located at GGSN's. All GGSN's may not have FA's. The SGSN and the GGSN may also be co-located. One FA in a PLMN is sufficient for offering MIP service, but for capacity and efficiency reasons, more than one may be desired. This means that the MS must request a PDP context to be set up with a GGSN that offers FA functionality. While setting up the PDP context, the MS is informed about network parameters of the FA, e.g. care-of address.

The MS may have the same care-of address COA, during a session, i.e. as long as a PDP context is activated. A very mobile MS might perform several inter-SGSN HO's during a long session which may cause an inefficient routing. As an initial improvement, a streamlining procedure, with a temporary anchoring point in the GGSN, could be introduced: If the MN is not transferring data, or possible even in the active state, while moving from one SGSN to another, a new PDP context could be setup between the new SGSN and its associated GGSN at the handover. The MN will get a new care-of address. If the MN is transferring data, e.g. being involved in a TCP session, the MN would move from the old SGSN to the new one while keeping the PDP Context in the old (anchor) GGSN for the duration of the data transfer. Once the data transfer is terminated, the PDP Context can be moved to the GGSN associated with the new SGSN and a new care-of address can be obtained.

The problem is how to discover the movement and to find a new foreign agent FA, preferably the nearest one, when the MN is moving from one SGSN to another. The GPRS terminal (MS) naturally is aware of the change of the SGSN on the GPRS protocol level as described above but this change is transparent to the overlaying MIP protocol and the mobile node MN associated to the GPRS terminal MS.

Disclosure of the Invention

An object of the present invention is to overcome or alleviate the above described problems.

The object is achieved by a method, a system and the access node which are characterized by what is disclosed in the attached independent

claims 1, 7 and 10, respectively. Preferred embodiments of the invention are disclosed in the attached dependent claims.

In the present invention a support node, or more generally any access node, which is a target of a handover during an IP session, is arranged to
5 check the optimal routing also in the mobile IP point of view. For that purpose, the access node is aware of the most preferred mobility agent, normally the closest one, which should be used. In a handover situation the system, preferably the access node, checks whether there is a more preferred mobility agent which should substitute for the current mobility agent of the IP session.
10 If there is no preferred mobility agent for that access node, or the preferred mobility agent appears to be the same as the current mobility agent of the IP session, the current mobility agent is maintained. However, if there is a more preferred mobility agent for that access node and the preferred mobility agent is not the same as the current mobility agent, the connection (e.g. a PDP
15 context) to the current mobility agent is closed (released), and a new connection (e.g. PDP context) to the preferred mobility agent of the respective access node is opened. A typical feature of the mobility agent is that it periodically transmits agent advertisement messages to the mobile nodes in order to advertise its services. The mobile nodes use these advertisements to determine
20 the current point of attachment to the Internet. In merit of the new connection established by the access node to the preferred mobility agent, the agent advertisement messages sent by the new mobility agent can be received by the mobile node, and thereby the mobile node is able to detect the change of the attachment point (i.e. mobility agent) and to initiate a standard mobile IP registration. Thus, an advantage of the invention is that the inventive new functionality at the access node enables to detect the movement of mobile node
25 also in the mobile IP level and to select and change the most optimal mobility agent in each part of the network, without needing any non-standard signalling or procedure in other elements of the packet radio network or on the mobile IP
30 level.

The route optimization saves the transmission resources in the packet radio system, and possibly also makes the connection faster as the connection leg between the access node and the mobility agent is shorter.

Brief description of the Drawings

In the following, the invention will be described in greater detail by means of preferred embodiments with reference to the accompanying drawings, in which

- 5 Figure 1 illustrates GPRS network architecture,
- Figure 2 is a signalling diagram illustrating the method according to the invention,
- Figure 3 is a flow diagram illustrating the function of the support node,
- 10 Figure 4 is a signalling diagram illustrating the method according to the invention.

Preferred Embodiments of the Invention

The present invention can be applied to any packet mode communications. The invention can be especially preferably used for providing a general packet radio service GPRS in the pan-European digital mobile communication system GSM (Global System for Mobile Communication) or in corresponding mobile communication systems, such as DCS1800 and PCS (Personal Communication System), or in third generation (3G) mobile systems, such as UMTS, implementing a GPRS-type packet radio. In the following, the preferred embodiments of the invention will be described by means of a GPRS packet radio network formed by the GPRS service and the 3G or GSM system without limiting the invention to this particular packet radio system.

25 A GPRS architecture utilizing a 3G radio access (such as UMTS) or a 2G radio access (such as GSM) is illustrated in Fig. 1. The GPRS infrastructure comprises support nodes such as a GPRS gateway support node (GGSN) and a GPRS serving support node (SGSN). The main functions of the GGSN nodes involve interaction with the external data network. The GGSN updates the location directory using routing information supplied by the SGSNs about an MS's path and routes the external data network protocol packet encapsulated over the GPRS backbone to the SGSN currently serving the MS. It also decapsulates and forwards external data network packets to the appropriate data network and handles the billing of data traffic.

35 The main functions of the SGSN are to detect new GPRS mobile stations in its service area, handle the process of registering the new MSs

along with the GPRS registers, send/receive data packets to/from the GPRS MS, and keep a record of the location of the MSs inside of its service area. The subscription information is stored in a GPRS register (HLR) where the mapping between a mobile's identity (such as MS-ISDN or IMSI) and the PSPDN address is stored. The GPRS register acts as a database from which the SGSNs can ask whether a new MS in its area is allowed to join the GPRS network.

The GPRS gateway support nodes GGSN connect an operator's GPRS network to external systems, such as other operators' GPRS systems, data networks 11, such as an IP network (Internet) or a X.25 network, and service centres. Fixed hosts 14 can be connected to data network 11 e.g. by means of a local area network LAN and a router 15. A border gateway BG provides an access to an inter-operator GPRS backbone network 12. The GGSN may also be connected directly to a private corporate network or a host. The GGSN includes GPRS subscribers' PDP addresses and routing information, i.e. SGSN addresses. Routing information is used for tunnelling protocol data units PDU from data network 11 to the current switching point of the MS, i.e. to the serving SGSN. The functionalities of the SGSN and GGSN can be connected to the same physical node (SGSN+GGSN).

The home location register HLR of the GSM network contains GPRS subscriber data and routing information and it maps the subscriber's IMSI into one or more pairs of the PDP type and PDP address. The HLR also maps each PDP type and PDP address pair into a GGSN node. The SGSN has a Gr interface to the HLR (a direct signalling connection or via an internal backbone network 13). The HLR of a roaming MS and its serving SGSN may be in different mobile communication networks.

The intra-operator backbone network 13, which interconnects an operator's SGSN and GGSN equipment can be implemented, for example, by means of a local network, such as an IP network. It should be noted that an operator's GPRS network can also be implemented without the intra-operator backbone network, e.g. by providing all features in one computer.

Network access is the means by which a user is connected to a telecommunication network in order to use the services and/or facilities of that network. An access protocol is a defined set of procedures that enables the user to employ the services and/or facilities of the network. The SGSN, which is at the same hierarchical level as the mobile switching centre MSC, keeps

track of the individual MSs' location and performs security functions and access control. GPRS security functionality is equivalent to the existing GSM security. The SGSN performs authentication and cipher setting procedures based on the same algorithms, keys, and criteria as in existing GSM. GPRS
5 uses a ciphering algorithm optimised for packet data transmission.

In order to access the GPRS services, a MS shall first make its presence known to the network by performing a GPRS attach. This operation establishes a logical link between the MS and the SGSN, and makes the MS available for SMS over GPRS, paging via SGSN, and notification of incoming
10 GPRS data. More particularly, when the MS attaches to the GPRS network, i.e. in a GPRS attach procedure, the SGSN creates a mobility management context (MM context), and a logical link LLC (Logical Link Control) established between the MS and the SGSN in a protocol layer. MM contexts, are stored in the SGSN and MS. The MM context of the SGSN may contain subscriber
15 data, such as the subscriber's IMSI, TLLI and location and routing information, etc.

In order to send and receive GPRS data, the MS shall activate the packet data address that it wants to use, by requesting a PDP activation procedure. This operation makes the MS known in the corresponding GGSN, and
20 interworking with external data networks can commence. More particularly, one or more PDP context is created in the MS and the GGSN and the SGSN, and stored in the serving SGSN in connection with the MM context. The PDP context defines different data transmission parameters, such as the PDP type (e.g. X.25 or IP), PDP address (e.g. IP address), quality of service QoS and
25 NSAPI (Network Service Access Point Identifier). The MS activates the PDU context with a specific message, Activate PDP Context Request, in which it gives information on the TLLI, PDP type, PDP address, required QoS and NSAPI, and optionally the access point name APN. The SGSN send a create PDP context message to the GGSN which creates the PDP context and send
30 it to SGSN. SGSN sends the PDP context to MS in a Activate PDP Context Response message, and a virtual connection or link between the MS and the GGSN is established. As a result the SGSN forwards all the data packets from the MS to the GGSN, and the GGSN forwards the SGSN all data packets received from the external network and addressed to the MS. The PDP context
35 is stored in the MS, the SGSN and the GGSN. When the MS roams to the

area of a new SGSN, the new SGSN requests MM and PDP contexts from the old SGSN.

Fig. 1 illustrates the implementation of mobile IP in the GPRS/3G environment.

5 The MS can be a laptop computer PC connected to a packet radio enabled cellular telephone. Alternatively, the MS can be an integrated combination of a small computer and a packet radio telephone, similar in appearance to the Nokia Communicator 9000 series. Yet further embodiments of the MS are various pagers, remote-control, surveillance and/or data-acquisition
10 devices, etc. The user of a mobile station MS subscribes to a special Mobile IP service. The subscription information is stored in the Home Location Register HLR together with the user's home IP address.

 In Fig. 1 the foreign agents FA are located at (integrated into) GGSN's. An alternative is that the SGSN and the GGSN are co-located, and
15 the FA's are located at SGSN+GGSNs. It should be noted that there may be more than one SGSN and GGSN in one network. All GGSN's may not have FA's. Each FA has an IP address in the Internet and in the operator's own private GPRS/3G backbone network. More precisely, the FA's IP address is such that IP packets destined to that address are routed in the Internet to the
20 GGSN associated with the FA. When the MN leaves its home subnet and registers to a new FA, it can no longer be reached on the basis of its home IP address alone, but must be assigned an address belonging to the visited network, called the *care-of address (COA)*. The care-of address positively identifies the instantaneous location of the mobile terminal and may be: 1) the IP
25 address of the FA belonging to the visited network, or 2) an IP address acquired directly by the mobile terminal through an autoconfiguration mechanism from the local IP address space, in which case the term co-located care-of address is used. When registering to a new FA and obtaining a COA, the MN which when registers with a home agent HA, in its home network, informing the latter of its COA. In Fig. 1 a home agent HA is located in a data network 11
30 which is the home network of the mobile node MN associated with the mobile station MS. A second host 14 wishing to communicate with the MN need not to be aware of the MN has moved: it simply sends IP packets addressed to MN's home IP address. These packets are routed via normal IP routing to the MN's home network, there they are intercepted by the HA. The HA encapsulates each such packet in another IP packet which contains the MN's COA as
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these packets are thus delivered to the FA (a process called tunneling). The FA forwards the IP packet to the GGSN. The GGSN forwards the IP packet (which may be encapsulated for transmission over the GPRS backbone) to the serving SGSN which further forwards the IP packet to the MS/MN. Packets from the MN to the another host 14 need not necessarily be tunneled: the MN may simply send them to the GGSN which directly forwards the packets to the second host 14, without interception by the FA or the HA.

As noted above, according to the present invention the SGSN determines whether it is preferable to change the mobility agent of the IP session or not. A preferred embodiment of the invention will be now described with reference to Figures 1, 2, 3 and 4.

A reference is now made to Figure 1. The home network of the mobile station MS is the GPRS/3G network 1. The user of the mobile station MS subscribes to a special mobile IP service, and an IP application in the MS or in a separate data terminal is a mobile node MN in a mobile IP communication. It is assumed that the MS/MN is attached to the home network 1 and the radio access network RAN1 (PS1 and PSC/RNC1). A serving support node in the home network is SGSN1. MM and PDP contexts have been created for mobile IP service as described above, and a virtual connection is provided between MS/MN and SGSN1 as well as between the SGSN1 and a gateway node GGSN1 which has an associated foreign agent FA1. Thus, the IP packets addressed to the MN can be forwarded to the MN over the home network 1 and RAN1. The COA of the MN has been registered to the home agent HA in the home network 11 of the MN, so that a mobile IP tunnelling is provided from the HA to the GGSN/FA1.

Let us now assume that the MS/MN moves to service area of another GPRS/3G network 2 which is served by a support node SGSN2. When the MS/MN arrives at a new RAN2, the MS part listens to radio broadcast messages, which contain information about radio parameters, network and cell identity, etc. as well as e.g. information about available core network, service providers, service capabilities etc. On the basis of the broadcast the MS determines that the network and/or the routing area has changed. Upon detecting a change of routing area, the MS/MN sends a routing area update request to the new SGSN, namely SGSN2, as shown in Figure 2. The new SGSN2 sends a SGSN context request message to the old SGSN1 (at step 2) to get the MN and PDP contexts for the MS/MN. The old SGSN1 responses with a

SGSN context response message which contain the MN and PDP contexts (step 3). At step 4 the new SGSN2 may, in certain situations, execute authentication/security functions which may involve an interrogation to the HLR of the MS/MN. If the user has at least one activated PDP context, when the new SGSN2 sends a SGSN context acknowledge message to the old SGSN1. The old SGSN1 may now start forwarding of buffered data packets belonging to the activated PDP context, if any, to the new SGSN2. The new SGSN2 will now execute the foreign agent check procedure according to the present invention, step 6. The FA check procedure according to the preferred embodiment of the present invention is illustrated in Figure 3. In step 31 the new SGSN2 checks whether there is a preferred FA defined for it. For example, SGSN2 may check whether there is an address of a preferred FA2 stored in the SGSN2. In this example, the address of GGSN/FA2 is found, and the procedure proceeds to step 32. In step 32 the new SGSN2 checks whether the address of the old FA1 obtained in the PDP context from the old SGSN1 is same as the stored address of the preferred FA2. In this example, the old FA1 is at GGSN1 and the preferred FA2 of the SGSN2 is at the GGSN2, and the addresses do not match. The procedure proceeds to the step 33 in which the new SGSN2 deletes the PDP context in the old GGSN/FA1 by sending a delete PDP context requests to the old GGSN/FA1, as shown in Figure 2. As a result any active PDP context in the GGSN/FA1 is deactivated, and the GGSN/FA1 acknowledges by sending a delete PDP context response to the new SGSN2 (step 8 in Figure 2). Referring again to Figure 3, the process proceeds to the step 34 wherein the new SGSN2 creates a PDP context in the preferred GGSN/FA2 by sending a create PDP context requests to the new GGSN/FA2 (step 9 in Figure 2). The GGSN/FA2 creates the PDP context for the MS/MN and returns a create PDP context response to the new SGSN2 (step 10 in Figure 2). The new SGSN2 establishes MN and PDP context for the MS/MN, and responses to the MS/MN with routing area update accept message (step 11). The MS/MN acknowledges with a routing area update complete message (step 12). A virtual connection has thus been established between the MS/MN and the GGSN/FA2.

All the previous procedures have been executed in the GPRS/3G layer only. The overlaying mobile IP layer and thereby the MN part of the MS/MN are not aware of a change of the FA. However, due to the newly established connection to the GGSN/FA2 the MN is able to hear the agent ad-

vertisement messages broadcasted by the new FA2 in accordance with the mobile IP protocol. Upon receiving the agent advertisement from the new FA2, the MN is able to detect a change in point of attachment, i.e. change of FA, in accordance with the MIP standard. The agent advertisement message may also include the care-of-address COA, or the MN may acquire the COA in accordance with the MIP standard. Then the mobile node MN registers its COA with its home agent HA in accordance with the MIP standard (step 14 in Figure 2). Depending on its method of attachment, the MN will register either directly with its HA, or through the new FA which forwards the registration to the HA.

10 Thereafter, the mobile IP tunnelling between the HA and the old GGSN/FA1 is released and the new mobile IP tunneling is established between the HA and the new GGSN/FA2, in accordance with the mobile procedures (step 15 in Figure 2).

As a results the change of FA have been detected and established using standard GPRS/3G procedures and messages and standard mobile IP procedures and messages everywhere else but in the SGSN2. Also in the SGSN2 only minor modifications are needed. Firstly, a preferred FA has to be defined for the SGSN2. Secondly, the need for FA change needs to be executed. Thirdly, the new SGSN is arranged to automatically and independently (without involvement of the MS) to delete the PDP context in the old GGSN and to create new PDP context in the new GGSN.

Referring again to Figure 3, if no preferred FA is defined for the SGSN1 in step 31, or the old FA is the same as the preferred FA (the addresses do match) in step 32, the process proceeds to step 35. In step 35 the new SGSN1 updates the PDP context in the old GGSN/FA1, as shown in step 41 in Figure 4. The old GGSN/FA1 updates the PDP context to contain the address of the new SGSN2, and sends an update PDP context response to the new SGSN2 in step 42. Then the new SGSN2 sends the RA update accept message to the MS/MN (step 43) and the MS/MN response with the RA update complete message (step 44). A virtual connection is thus established between the MS/MN and the old GGSN/FA via the new SGSN1. As the FA and the COA are unchanged, no registration to the HA is needed. In Figure 4 steps 1-6 are similar to those in Figure 2.

The description only illustrates preferred embodiments of the invention. The invention is not, however, limited to these examples, but it may vary within the scope and spirit of the appended claims.

Claims

1. A method of providing Internet Protocol-type, or IP-type, mobility for a mobile node in a packet radio system comprising a plurality of mobile nodes (MS/MN), a first and a second access node (SGN1, SGN2) serving said mobile nodes within first (RAN1) and second (RAN2) parts of the packet radio system, respectively, at least one first gateway node (GGSN1) for interfacing said first part (RAN1) of the packet radio system with external networks (12), and a first mobility agent (FA1) which is associated with said at least one first gateway node (GGSN1) and arranged to provide mobile IP routing services to the mobile nodes (MS/MN) while registered to the first part (RAN1) of the packet radio system, said method comprising steps of

establishing an IP session between one of said plurality of mobile nodes (MS/MN) and a second party via said first access node (SGSN1) and said first mobility agent (FA1),

rerouting the IP session via said second access node (SGSN2) in response to a movement of said one of mobile nodes (MS(MN) to said second part (RAN2) of the packet radio system,

characterized by further steps of

checking whether there is a second mobility agent (FA2) which is more preferred in a routing point of view than said first mobility agent (FA1), and

reacting to said checking by

A) maintaining a connection to said first mobility agent (FA1), if there is no second mobility agent which is more preferred than said first one, and

B) closing the connection to said first mobility agent (FA1) and opening new connection to said second mobility agent (FA2), if said more preferred second mobility agent is available, and sending an agent advertisement message from said second mobility agent (FA2) to said one mobile node (MS/MN) over said new connection, said agent advertising message enabling said one mobile node (MS/MN) to detect a change of attachment point and to initiate a mobile IP registration.

2. The method according to claim 1, characterized by steps of

storing in said second access node (SGSN2) an identity of said preferred mobility agent (FA2) of said second access node,

checking at said second access node (SGSN2), in response to a movement of said one mobile node from said first access node (SGSN1) to said second access node (SGSN2), whether an identity of said first mobility agent (FA1) and said stored identity of said preferred mobility agent (FA2) do
 5 match or not,

maintaining a connection to said first mobility agent (FA1), if the identities do match, and

closing the connection to said first mobility agent (FA1) and opening new connection to said preferred mobility agent (FA2), if said identities no not
 10 match.

3. The method according to claim 1 or 2, characterized in that said steps of closing and opening of the connection comprises steps of closing and opening, respectively, a packet radio protocol context .

4. The method according to claim 1, 2 or 3, characterized in
 15 that said preferred mobility agent of said second access node is a foreign agent (FA2) associated with a gateway node (GGSN2) in said second part (RAN2) of the packet radio network.

5. The method according to claim 2, 3 or 4, , characterized in that said identity includes an address of mobility agent (FA1, FA2).

20 6. A packet radio system, comprising
 a plurality of mobile nodes (MS/MN),
 a first (SGSN1) and a second (SGSN2) access node serving said mobile nodes within first (RAN1) and second (RAN2) parts of the packet radio system, respectively,

25 at least one first gateway node (GGSN1) for interfacing said first part of the packet radio system with external networks (11),

a first mobility agent (FA1) which is associated with said at least one first gateway node (GGSN1) and arranged to route a IP connection to any one of said the mobile nodes via said while registered to the first part (RAN1)
 30 of the packet radio system,

a rerouting mechanism by which said IP connection initially routed via said first access node (SGSN1) and said first mobility agent (FA1) can be routed via said second access node (SGSN2) in response to a movement of said one of mobile nodes (MS/MN) to said second part (RAN2) of the packet
 35 radio system,

characterized by said rerouting mechanism comprising

the system being arranged to check whether there is a second mobility agent (FA2) which is more preferred in a routing point of view than said first mobility agent (FA1),

the system being arranged to close a connection to said first mobility agent (FA1) and opening a new connection to said second mobility agent (FA2), if said more preferred second mobility agent (FA2) is available according to said checking,

said second mobility agent (FA2) being arranged to send an agent advertisement message to said one mobile node (MS/MN) over said new connection,

said mobile node (MS/MN) being arranged to detect a change of attachment by means of said agent advertising message and to initiate a mobile IP registration.

7. The system according to claim 6, characterized in that said preferred mobility agent of said second access node is a foreign agent (FA2) associated with a gateway node (GGSN2) in said second part (RAN2) of the packet radio network.

8. The system according to claim 6 or 7, characterized in that said second access node (SGSN2) is arranged to make said checking.

9. The system according to claim 6, 7 or 8, characterized in that the second access node (SGSN2) is arranged to store an identity, such as address, of the preferred foreign agent (FA2) for the checking purposes.

10. An access node for a packet radio system comprising a plurality of mobile nodes (MS/MN), access nodes (SGSN1, SGSN2) serving said mobile nodes within respective parts (RAN1, RAN2) of the packet radio system, at least two gateway (GGSN1, GGSN2) nodes for interfacing the packet radio system with external networks (11), and at least two mobility agents (FA1, FA2) which are associated with different ones of said at least two gateway nodes (GGSN1, GGSN2) and arranged to provide mobile IP routing services to the mobile nodes (MS/MN) while registered to the packet radio system, characterized by said access node comprising

means for checking, when a mobile node (MS/MN) having an IP connection through another access node (SGSN1) and a first mobility agent (FA1) is accessing the system via said access node (SGSN2), whether there is another mobility agent which is more preferred in a routing point of view than said first mobility agent (FA1),

means responsive to said checking means for closing a connection to said first mobility agent (FA1) and opening a new connection to said preferred other mobility agent (FA2), if said more preferred other mobility agent (FA2) is available.

5 11. The access node according to claim 10, characterized in that said checking means comprises

means for storing an identity, such as address, of said preferred other mobility agent (FA2) of said access node (SGSN2),

10 means for checking, in response to a movement of said mobile node (MS/MN) from said other access node (SGSN1) to said access node (SGSN2), whether an identity of said first mobility agent (FA1) and said stored identity of said preferred mobility agent (FA2) do match or not.

12. The access node according to claim 10 or 11, characterized in that said closing and opening means comprise

15 means for maintaining a connection to said first mobility agent (FA1), if the identities do match, and

means for closing the connection to said first mobility agent (FA1) and opening new connection to said preferred mobility agent (FA2), if said identities no not match.

20 13. The access node according to claim 10, 11 or 12, characterized in that said means for closing and opening of the connection comprise means for closing a packet radio protocol context in the gateway node (GGSN1) of the first mobility agent (FA1) and opening a packet radio protocol context in the gateway node (GGSN2) of the preferred mobility agent (FA2).

25 14. The access node according to claim 13, characterized in that said means for maintaining the connection comprise means for updating a packet radio protocol context of the mobile node (MS/MN) in the gateway node (GGSN1) of the first mobility agent (FA1).

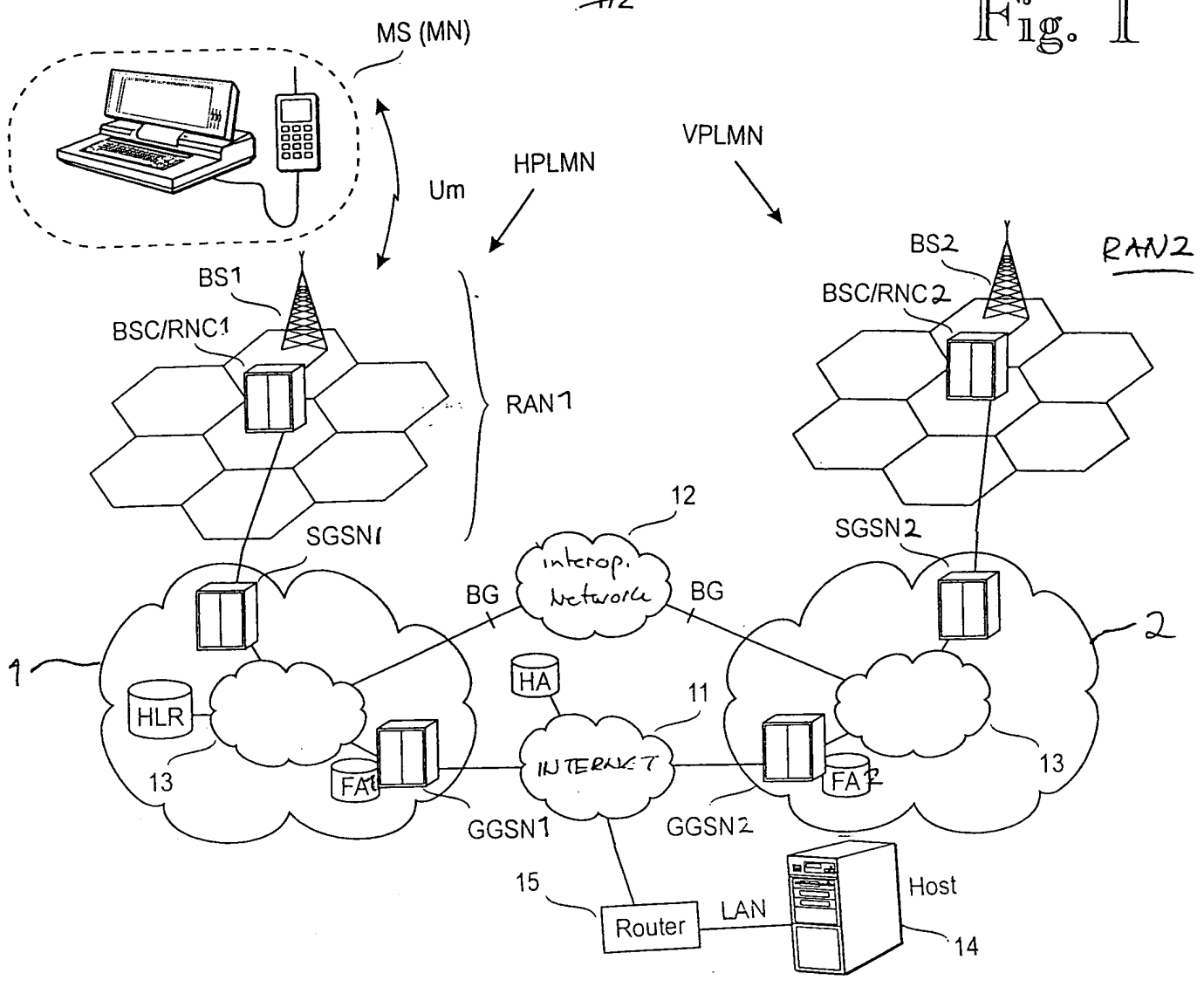
(57) Abstract

In a packet radio network which supports a mobile IP protocol, A very mobile MS might perform several handovers between access nodes during a long session which may cause an inefficient mobile IP routing. In the present invention, an access node which is a target of a handover during an IP session, is arranged to check the optimal routing also in the mobile IP point of view. For that purpose, the access node is aware of the most preferred mobility agent, normally the closest one, which should be used. In a handover situation the system, preferably the access node, checks (31,32) whether there is a more preferred mobile IP mobility agent which should substitute for the current mobility agent of the IP session. If there is no preferred mobility agent for that access node, or the preferred mobility agent appears to be the same as the current mobility agent of the IP session, the current mobility agent is maintained (35). However, if there is a more preferred mobility agent for that access node and the preferred mobility agent is not the same as the current mobility agent, the connection the current mobility agent is closed (33), and a new connection to the preferred mobility agent of the respective access node is opened (34). In merit of the new connection established by the access node to the preferred mobility agent, the agent advertisement messages sent by the new mobility agent can be received by the mobile node, and thereby the mobile node is able to detect the change of the attachment point (i.e. mobility agent) and to initiate a standard mobile IP registration.

(Fig. 3)

2 3
1/2

Fig. 1



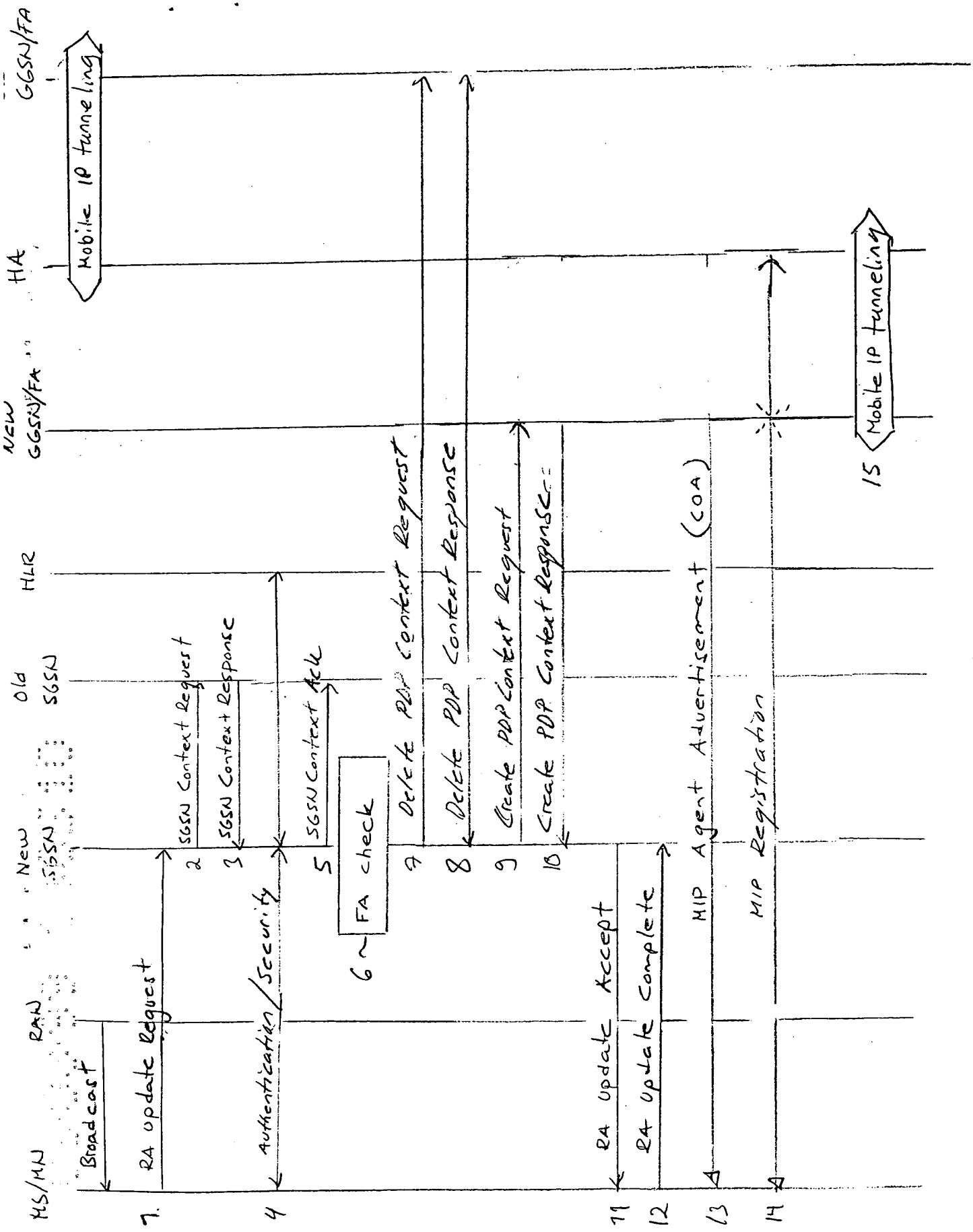


Fig. 2.

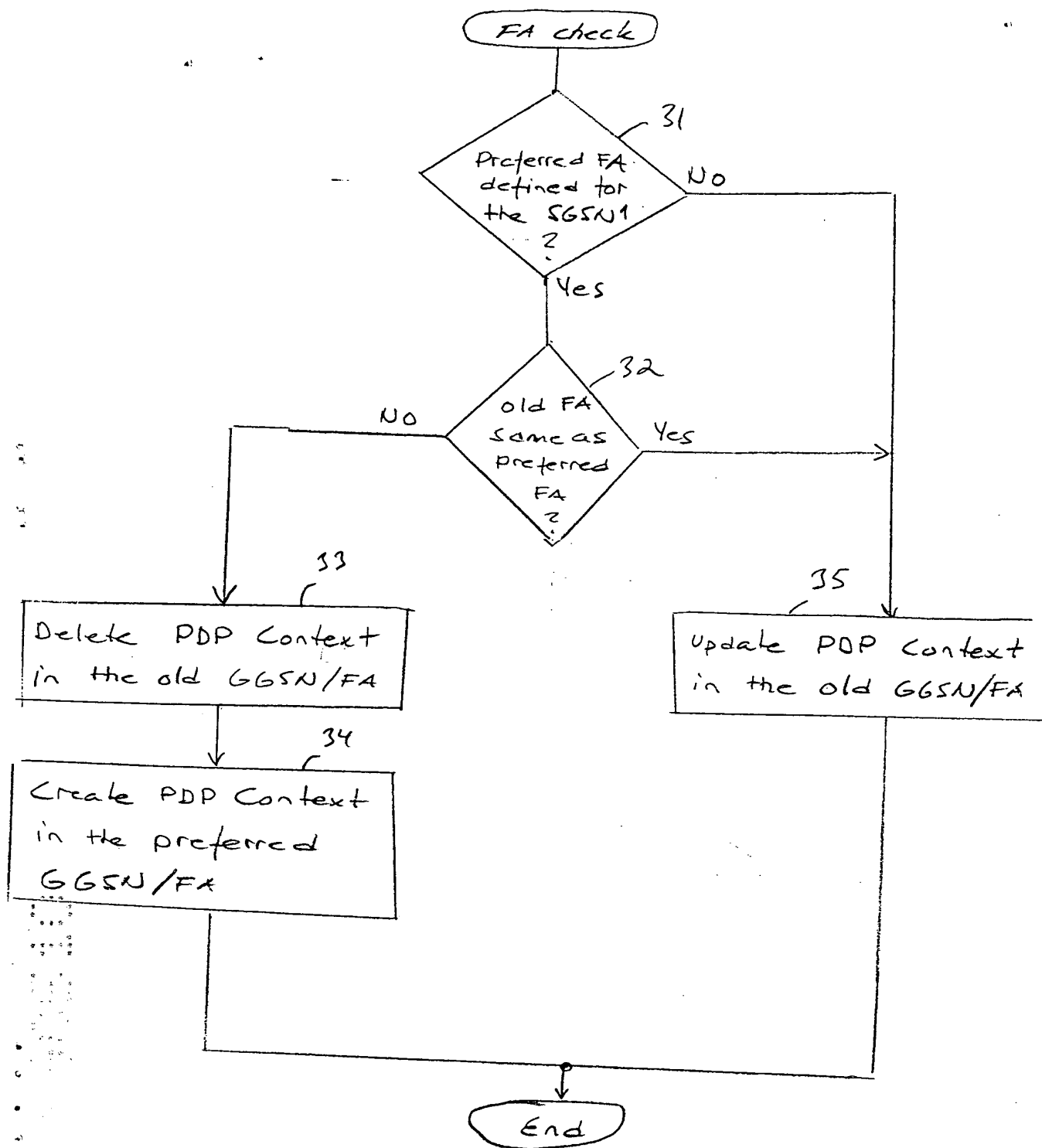


Fig. 3

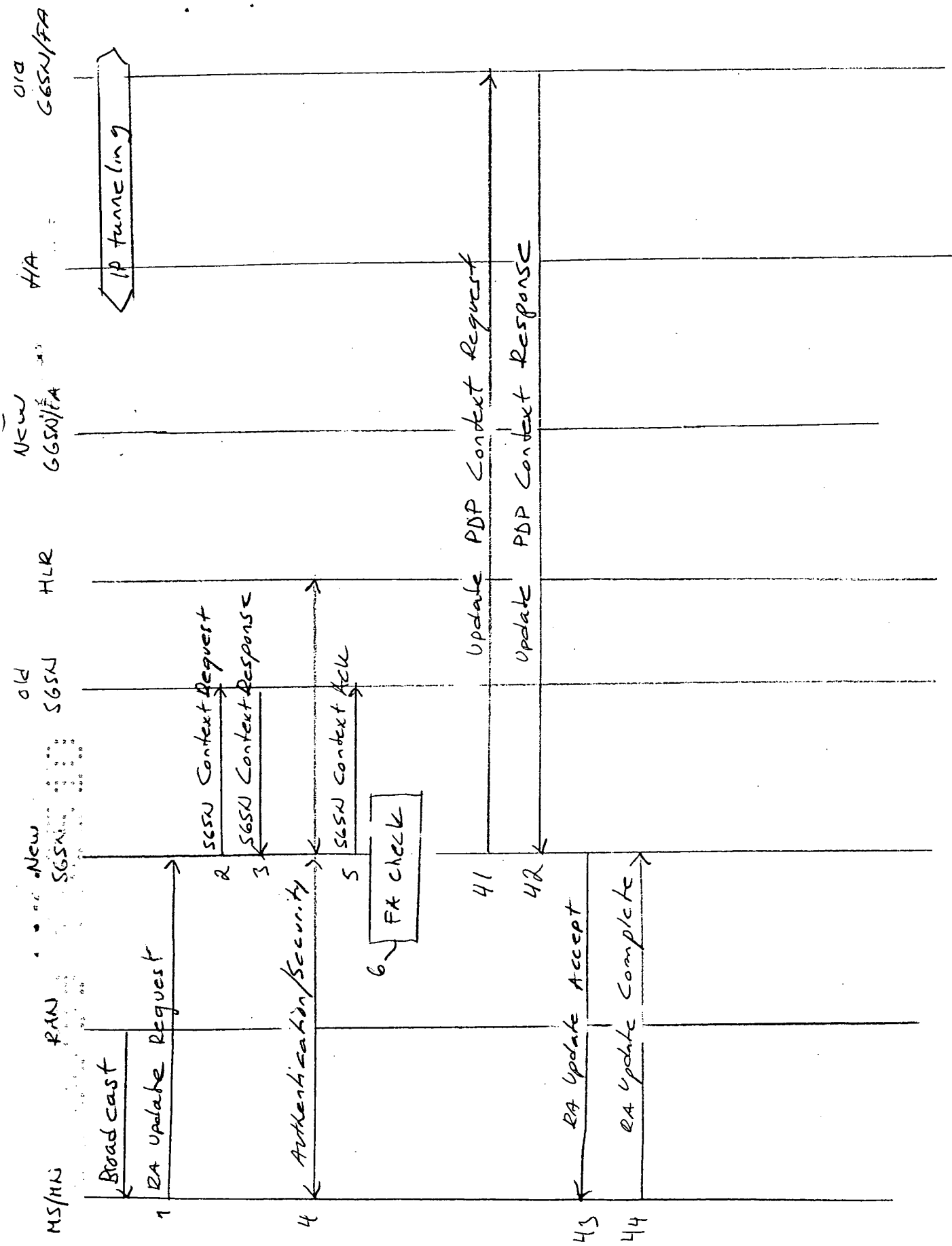


Fig. 4